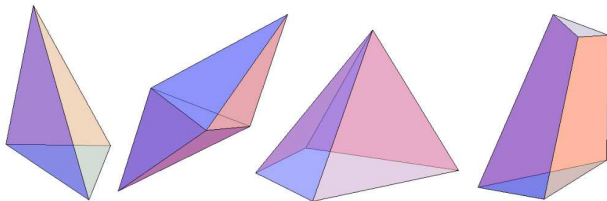


Barcelona - January 26, 2018



Slack ideals of Polytopes

[arXiv:1708.04739]

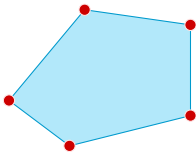
Antonio Macchia

joint work with **João Gouveia** (U Coimbra)
Rekha Thomas, Amy Wiebe (U Washington)

Polytopes

A ***polytope*** P is a convex hull of a finite set of points in some \mathbb{R}^d :

$$P = \text{conv}(p_1, \dots, p_v)$$

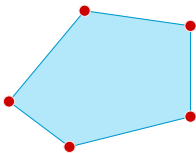


\mathcal{V} -*representation*

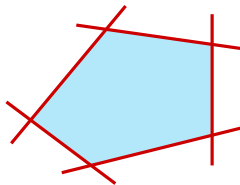
Polytopes

A **polytope** P is a convex hull of a finite set of points in some \mathbb{R}^d :

$$P = \text{conv}(p_1, \dots, p_v) = \{\mathbf{x} \in \mathbb{R}^d : A\mathbf{x} \leq \mathbf{b}\}.$$



\mathcal{V} -representation

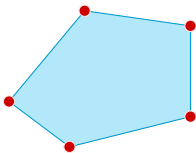


\mathcal{H} -representation

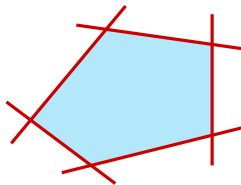
Polytopes

A **polytope** P is a convex hull of a finite set of points in some \mathbb{R}^d :

$$P = \text{conv}(p_1, \dots, p_v) = \{\mathbf{x} \in \mathbb{R}^d : A\mathbf{x} \leq \mathbf{b}\}.$$



\mathcal{V} -representation

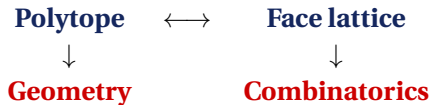


\mathcal{H} -representation

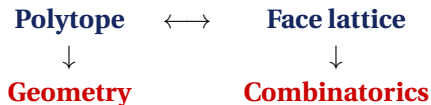
The **dimension** of a polytope is the dimension of its affine hull.

The *face lattice* of P is the set of all faces of P , ordered by inclusion.

The *face lattice* of P is the set of all faces of P , ordered by inclusion.



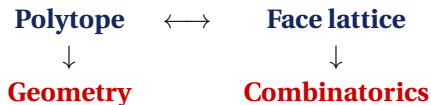
The *face lattice* of P is the set of all faces of P , ordered by inclusion.



Find all polytopes of a fixed dimension:

- combinatorial types of polytopes \leftrightarrow finite lattices corresponding to face lattices of polytopes,
- describe the set of all realizations of a given combinatorial type (realization space).

The *face lattice* of P is the set of all faces of P , ordered by inclusion.



Find all polytopes of a fixed dimension:

- combinatorial types of polytopes \leftrightarrow finite lattices corresponding to face lattices of polytopes,
- describe the set of all realizations of a given combinatorial type (realization space).

Theorem (Steinitz, 1922) *A graph G is the edge graph of a 3-polytope $\Leftrightarrow G$ is simple, planar and 3-connected.*

Realization Space

Given: combinatorial type of a d -polytope

→ face lattice or vertex-facet incidences

Want: all ways to realize this type in \mathbb{R}^d

Realization Space

Given: combinatorial type of a d -polytope

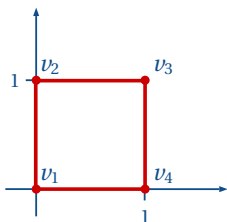
→ face lattice or vertex-facet incidences

Want: all ways to realize this type in \mathbb{R}^d

P is a quadrilateral

4 vertices $\{v_1, v_2, v_3, v_4\}$

4 facets $\{v_1, v_2\}, \{v_2, v_3\}, \{v_3, v_4\}, \{v_4, v_1\}$



Combinatorial equivalence

$Q \stackrel{c}{=} P \Leftrightarrow P$ and Q have the same vertex-facet incidences

All quadrilaterals are combinatorially equivalent to a square

Realization Space

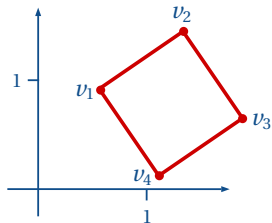
Given: combinatorial type of a d -polytope
→ face lattice or vertex-facet incidences

Want: all ways to realize this type in \mathbb{R}^d

P is a quadrilateral

4 vertices $\{v_1, v_2, v_3, v_4\}$

4 facets $\{v_1, v_2\}, \{v_2, v_3\}, \{v_3, v_4\}, \{v_4, v_1\}$



Affine equivalence

$$Q \stackrel{a}{=} P \Leftrightarrow Q = \psi(P), \quad \psi(\mathbf{x}) = A\mathbf{x} + \mathbf{b}$$

preserves parallel lines, e.g. scaling, rotation, reflection, translation

Parallelograms are affinely equivalent to a square

Realization Space

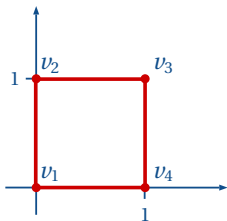
Given: combinatorial type of a d -polytope
→ face lattice or vertex-facet incidences

Want: all ways to realize this type in \mathbb{R}^d

P is a quadrilateral

4 vertices $\{v_1, v_2, v_3, v_4\}$

4 facets $\{v_1, v_2\}, \{v_2, v_3\}, \{v_3, v_4\}, \{v_4, v_1\}$



Projective equivalence

$$Q \stackrel{p}{=} P \Leftrightarrow Q = \phi(P), \quad \phi(\mathbf{x}) = \frac{A\mathbf{x} + \mathbf{b}}{\mathbf{c}^\top \mathbf{x} + d}, \quad \det \begin{bmatrix} A & \mathbf{b} \\ \mathbf{c}^\top & d \end{bmatrix} \neq 0$$

All convex quadrilaterals are projectively equivalent to a square.

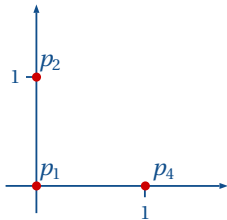
A square is **projectively unique**.

Realization Space

Set of all realizations of polytopes combinatorially equivalent to P

Mod out affine transformation:

fix an **affine** basis B of $d + 1$ common vertices

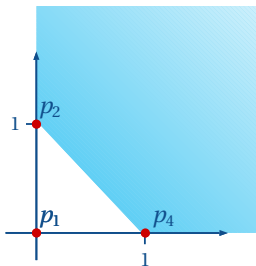


Realization Space

Set of all realizations of polytopes combinatorially equivalent to P

Mod out affine transformation:

fix an **affine** basis B of $d + 1$ common vertices

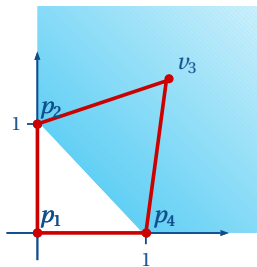


Realization Space

Set of all realizations of polytopes combinatorially equivalent to P

Mod out affine transformation:

fix an **affine** basis B of $d + 1$ common vertices



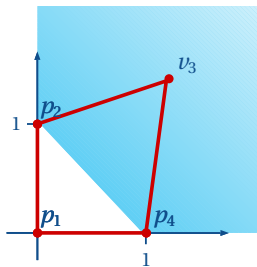
$$v_3 \in \left\{ (x, y) : \begin{array}{l} x > 0, y > 0 \\ x + y < 1 \end{array} \right\}$$

Realization Space

Set of all realizations of polytopes combinatorially equivalent to P

Mod out affine transformation:

fix an **affine** basis B of $d + 1$ common vertices



$$v_3 \in \left\{ (x, y) : \begin{array}{l} x > 0, y > 0 \\ x + y < 1 \end{array} \right\}$$

Realization space

$$\mathcal{R}(P, B) = \left\{ Q = \text{conv}\{q_1, \dots, q_n\} \subset \mathbb{R}^d : q_i = p_i \forall p_i \in B, Q \stackrel{c}{=} P \right\}$$

Main Results

- A model for the realization space of a polytope up to projective equivalence, that arises as the positive part of an algebraic variety.
- Naturally mods out affine equivalence, so no choice of basis is needed.
- The ideal defining the variety is a computational engine for questions about realizations.
- The ideal suggests a new way to classify polytopes.

Slack Matrices

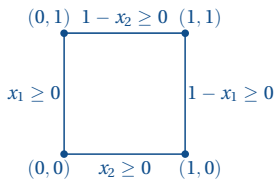
P : d -polytope in \mathbb{R}^d

vertices: $\{p_1, \dots, p_v\}$

facet inequalities: $\beta_1 - a_1^T \mathbf{x} \geq 0$

\vdots

$\beta_f - a_f^T \mathbf{x} \geq 0$



Slack Matrices

P : d -polytope in \mathbb{R}^d

vertices: $\{p_1, \dots, p_v\}$

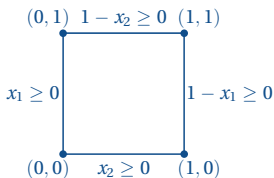
facet inequalities: $\beta_1 - a_1^\top x \geq 0$

\vdots

$\beta_f - a_f^\top x \geq 0$

Slack matrix

$$S_P = \begin{pmatrix} \vdots & & & \\ \cdots & \beta_j - a_j^\top p_i & \cdots & \\ \vdots & & & \end{pmatrix}_{v \times f}$$



$$S_P = \begin{pmatrix} 0 & 1 & 1 & 0 \\ 0 & 0 & 1 & 1 \\ 1 & 0 & 0 & 1 \\ 1 & 1 & 0 & 0 \end{pmatrix}$$

Slack Matrices

P : d -polytope in \mathbb{R}^d

vertices: $\{p_1, \dots, p_v\}$

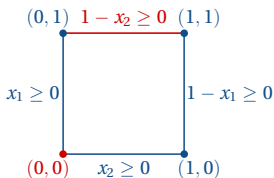
facet inequalities: $\beta_1 - a_1^T \mathbf{x} \geq 0$

\vdots

$\beta_f - a_f^T \mathbf{x} \geq 0$

Slack matrix

$$S_P = \begin{pmatrix} \vdots & & & \\ \cdots & \beta_j - a_j^T p_i & \cdots & \\ \vdots & & & \end{pmatrix}_{v \times f}$$



$$S_P = \begin{pmatrix} 0 & \mathbf{1} & 1 & 0 \\ 0 & 0 & 1 & 1 \\ 1 & 0 & 0 & 1 \\ 1 & 1 & 0 & 0 \end{pmatrix}$$

Slack Matrices

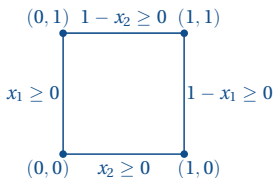
P : d -polytope in \mathbb{R}^d

vertices: $\{p_1, \dots, p_v\}$

facet inequalities: $\beta_1 - a_1^T x \geq 0$

\vdots

$\beta_f - a_f^T x \geq 0$



Slack matrix

$$S_P = \begin{pmatrix} \vdots & & & \\ \cdots & \beta_j - a_j^T p_i & \cdots & \\ \vdots & & & \end{pmatrix}_{v \times f}$$

$$S_P = \begin{pmatrix} 0 & 1 & 1 & 0 \\ 0 & 0 & 1 & 1 \\ 1 & 0 & 0 & 1 \\ 1 & 1 & 0 & 0 \end{pmatrix}$$

zero pattern \leftrightarrow combinatorics

$\text{rank}(S_P) = d + 1$

Slack Matrices

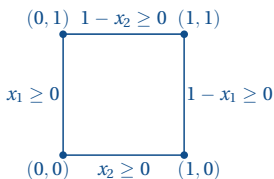
P : d -polytope in \mathbb{R}^d

vertices: $\{p_1, \dots, p_v\}$

facet inequalities: $\beta_1 - a_1^T x \geq 0$

\vdots

$\beta_f - a_f^T x \geq 0$



Slack matrix

$$S_P = \begin{pmatrix} \vdots & & & \\ \cdots & \beta_j - a_j^T p_i & \cdots & \\ \vdots & & & \end{pmatrix}_{v \times f}$$

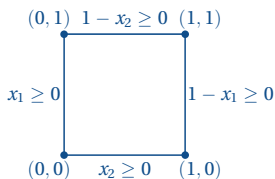
$$S_P = \begin{pmatrix} 0 & 1 & 1 & 0 \\ 0 & 0 & 1 & 1 \\ 1 & 0 & 0 & 1 \\ 1 & 1 & 0 & 0 \end{pmatrix}$$

zero pattern \leftrightarrow combinatorics $\text{rank}(S_P) = d + 1$

- P has infinitely many slack matrices $\{S_P D_f : D_f \text{ pos. diag. matrix}\}$

Slack Matrices

- P : d -polytope in \mathbb{R}^d
vertices: $\{p_1, \dots, p_v\}$
facet inequalities: $\beta_1 - a_1^\top x \geq 0$
 \vdots
 $\beta_f - a_f^\top x \geq 0$



Slack matrix

$$S_P = \begin{pmatrix} \vdots & & & \\ \cdots & \beta_j - a_j^\top p_i & \cdots & \\ \vdots & & & \end{pmatrix}_{v \times f} \quad S_P = \begin{pmatrix} 0 & 1 & 1 & 0 \\ 0 & 0 & 1 & 1 \\ 1 & 0 & 0 & 1 \\ 1 & 1 & 0 & 0 \end{pmatrix}$$

zero pattern \leftrightarrow combinatorics $\quad \text{rank}(S_P) = d + 1$

- P has infinitely many slack matrices $\{S_P D_f : D_f \text{ pos. diag. matrix}\}$

Affine equivalence

- $P \stackrel{a}{=} Q \Leftrightarrow P$ and Q have the same slack matrices

Lemma (GGKPRT, 2013) S slack matrix of $P \Rightarrow \text{conv}(\text{rows}(S))$ is affinely equivalent to P .

Combinatorial equivalence

$$S_P = \left(\begin{array}{ccc} & \vdots & \\ \cdots & \beta_j & -\mathbf{a}_j^\top \mathbf{p}_i \cdots \\ & \vdots & \end{array} \right)_{v \times f} = \underbrace{\left(\begin{array}{cc} 1 & \mathbf{p}_1^\top \\ \vdots & \vdots \\ 1 & \mathbf{p}_v^\top \end{array} \right)}_{v \times (d+1)} \underbrace{\left(\begin{array}{ccc} \beta_1 & \cdots & \beta_f \\ -\mathbf{a}_1 & \cdots & -\mathbf{a}_f \end{array} \right)}_{(d+1) \times f}$$

Combinatorial equivalence

$$S_P = \left(\begin{array}{ccc} & \vdots & \\ \cdots & \beta_j & -\mathbf{a}_j^\top \mathbf{p}_i \cdots \\ & \vdots & \end{array} \right)_{v \times f} = \underbrace{\left(\begin{array}{c} 1 \\ \vdots \\ 1 \end{array} \right)}_{v \times (d+1)} \underbrace{\left(\begin{array}{ccc} \mathbf{p}_1^\top & \cdots & \\ -\mathbf{a}_1 & \cdots & -\mathbf{a}_f \end{array} \right)}_{(d+1) \times f}$$

Theorem (GGKPRT, 2013) *A nonnegative matrix S is the slack matrix of some realization of P if and only if*

- 1 $\text{supp}(S) = \text{supp}(S_P)$;
- 2 $\text{rank}(S) = \text{rank}(S_P) = d + 1$;
- 3 *the all ones vector lies in the column span of S .*

Combinatorial equivalence

$$S_P = \left(\begin{array}{ccc} & \vdots & \\ \cdots & \beta_j - \mathbf{a}_j^\top \mathbf{p}_i & \cdots \\ & \vdots & \end{array} \right)_{v \times f} = \underbrace{\left(\begin{array}{c} 1 \\ \vdots \\ 1 \end{array} \right)_{v \times (d+1)}}_{v \times (d+1)} \underbrace{\left(\begin{array}{ccc} \mathbf{p}_1^\top & & \\ & \ddots & \\ & & \mathbf{p}_v^\top \end{array} \right)}_{(d+1) \times f} \underbrace{\left(\begin{array}{ccc} \beta_1 & \cdots & \beta_f \\ -\mathbf{a}_1 & \cdots & -\mathbf{a}_f \end{array} \right)}_{(d+1) \times f}$$

Theorem (GGKPRT, 2013) A nonnegative matrix S is the slack matrix of some realization of P if and only if

- 1 $\text{supp}(S) = \text{supp}(S_P)$;
- 2 $\text{rank}(S) = \text{rank}(S_P) = d + 1$;
- 3 the all ones vector lies in the column span of S .

Projective equivalence

Theorem (GPRT, 2017)

$Q \stackrel{P}{=} P \Leftrightarrow S_Q = D_v S_P D_f$ for some positive diagonal matrices D_v, D_f

Slack Ideal

Symbolic slack matrix

Replace nonzero entries of S_P by distinct variables.

$$S_P = \begin{pmatrix} 0 & 1 & 1 & 0 \\ 0 & 0 & 1 & 1 \\ 1 & 0 & 0 & 1 \\ 1 & 1 & 0 & 0 \end{pmatrix} \rightarrow S_P(\mathbf{x}) = \begin{pmatrix} 0 & x_1 & x_2 & 0 \\ 0 & 0 & x_3 & x_4 \\ x_5 & 0 & 0 & x_6 \\ x_7 & x_8 & 0 & 0 \end{pmatrix}$$

Slack Ideal

Symbolic slack matrix

Replace nonzero entries of S_P by distinct variables.

$$S_P = \begin{pmatrix} 0 & 1 & 1 & 0 \\ 0 & 0 & 1 & 1 \\ 1 & 0 & 0 & 1 \\ 1 & 1 & 0 & 0 \end{pmatrix} \rightarrow S_P(\mathbf{x}) = \begin{pmatrix} 0 & x_1 & x_2 & 0 \\ 0 & 0 & x_3 & x_4 \\ x_5 & 0 & 0 & x_6 \\ x_7 & x_8 & 0 & 0 \end{pmatrix}$$

Slack ideal

$$I_P = \langle (d+2)\text{-minors of } S_P(\mathbf{x}) \rangle : \left(\prod x_i \right)^\infty$$

$$I_P = \langle x_2 x_4 x_5 x_8 - x_1 x_3 x_6 x_7 \rangle$$

Slack realization space

- Positive part of slack variety: $\mathcal{V}_+(I_P) = \mathcal{V}(I_P) \cap \mathbb{R}_+^n$

Slack realization space

- Positive part of slack variety: $\mathcal{V}_+(I_P) = \mathcal{V}(I_P) \cap \mathbb{R}_+^n$
- $\mathbb{R}_{>0}^v \times \mathbb{R}_{>0}^f$ acts on $\mathcal{V}_+(I_P)$:

$$D_v \mathbf{s} D_f \in \mathcal{V}_+(I_P) \quad \text{for every } \mathbf{s} \in \mathcal{V}_+(I_P),$$

D_v, D_f positive diagonal matrices

Slack realization space

- Positive part of slack variety: $\mathcal{V}_+(I_P) = \mathcal{V}(I_P) \cap \mathbb{R}_+^n$
- $\mathbb{R}_{>0}^v \times \mathbb{R}_{>0}^f$ acts on $\mathcal{V}_+(I_P)$:

$$D_v \mathbf{s} D_f \in \mathcal{V}_+(I_P) \quad \text{for every } \mathbf{s} \in \mathcal{V}_+(I_P),$$

D_v, D_f positive diagonal matrices

Theorem (Gouveia, M, Thomas, Wiebe, 2017)

- 1 $\mathcal{V}_+(I_P) = \{ \text{nonneg. } S \text{ with } \text{supp}(S) = \text{supp}(S_P), \text{rank}(S) = d + 1 \}.$

Slack realization space

- Positive part of slack variety: $\mathcal{V}_+(I_P) = \mathcal{V}(I_P) \cap \mathbb{R}_+^n$
- $\mathbb{R}_{>0}^v \times \mathbb{R}_{>0}^f$ acts on $\mathcal{V}_+(I_P)$:

$$D_v \mathbf{s} D_f \in \mathcal{V}_+(I_P) \quad \text{for every } \mathbf{s} \in \mathcal{V}_+(I_P),$$

D_v, D_f positive diagonal matrices

Theorem (Gouveia, M, Thomas, Wiebe, 2017)

- 1 $\mathcal{V}_+(I_P) = \{ \text{nonneg. } S \text{ with } \text{supp}(S) = \text{supp}(S_P), \text{rank}(S) = d + 1 \}$.
- 2 $\mathcal{V}_+(I_P) / (\mathbb{R}_{>0}^v \times \mathbb{R}_{>0}^f) \xrightarrow{1:1}$ classes of projectively equivalent polytopes of the same combinatorial type as P .

Slack realization space

- Positive part of slack variety: $\mathcal{V}_+(I_P) = \mathcal{V}(I_P) \cap \mathbb{R}_+^n$
- $\mathbb{R}_{>0}^v \times \mathbb{R}_{>0}^f$ acts on $\mathcal{V}_+(I_P)$:

$$D_v \mathbf{s} D_f \in \mathcal{V}_+(I_P) \quad \text{for every } \mathbf{s} \in \mathcal{V}_+(I_P), \\ D_v, D_f \text{ positive diagonal matrices}$$

Theorem (Gouveia, M, Thomas, Wiebe, 2017)

- 1 $\mathcal{V}_+(I_P) = \{ \text{nonneg. } S \text{ with } \text{supp}(S) = \text{supp}(S_P), \text{rank}(S) = d + 1 \}$.
- 2 $\mathcal{V}_+(I_P) / (\mathbb{R}_{>0}^v \times \mathbb{R}_{>0}^f) \xleftrightarrow{1:1}$ classes of projectively equivalent polytopes of the same combinatorial type as P .

We call $\mathcal{V}_+(I_P) / (\mathbb{R}_{>0}^v \times \mathbb{R}_{>0}^f)$ the **slack realization space** of P .

Application 1: Realizability

Steinitz problem *Check whether an abstract polytopal complex is the boundary of an actual polytope.*

Application 1: Realizability

Steinitz problem *Check whether an abstract polytopal complex is the boundary of an actual polytope.*

[Altshuler, Steinberg, 1985]: 4-polytopes and 3-spheres with 8 vertices.

The smallest non-polytopal 3-sphere has vertex-facet non-incidence matrix

$$S_P(\mathbf{x}) = \begin{pmatrix} 0 & 0 & 0 & 0 & 0 & x_1 & x_2 & x_3 & x_4 & x_5 \\ 0 & 0 & 0 & 0 & x_6 & x_7 & 0 & 0 & x_8 & x_9 \\ 0 & 0 & x_{10} & x_{11} & x_{12} & 0 & 0 & 0 & 0 & x_{13} \\ 0 & 0 & x_{14} & x_{15} & 0 & 0 & x_{16} & x_{17} & 0 & 0 \\ 0 & x_{18} & 0 & x_{19} & 0 & 0 & 0 & x_{20} & x_{21} & x_{22} \\ x_{23} & 0 & x_{24} & 0 & 0 & x_{25} & x_{26} & 0 & 0 & 0 \\ x_{27} & x_{28} & 0 & 0 & x_{29} & 0 & 0 & 0 & 0 & 0 \\ x_{30} & x_{31} & 0 & 0 & 0 & 0 & x_{32} & x_{33} & x_{34} & 0 \end{pmatrix}.$$

Application 1: Realizability

Steinitz problem *Check whether an abstract polytopal complex is the boundary of an actual polytope.*

[Altshuler, Steinberg, 1985]: 4-polytopes and 3-spheres with 8 vertices.

The smallest non-polytopal 3-sphere has vertex-facet non-incidence matrix

$$S_P(\mathbf{x}) = \begin{pmatrix} 0 & 0 & 0 & 0 & 0 & x_1 & x_2 & x_3 & x_4 & x_5 \\ 0 & 0 & 0 & 0 & x_6 & x_7 & 0 & 0 & x_8 & x_9 \\ 0 & 0 & x_{10} & x_{11} & x_{12} & 0 & 0 & 0 & 0 & x_{13} \\ 0 & 0 & x_{14} & x_{15} & 0 & 0 & x_{16} & x_{17} & 0 & 0 \\ 0 & x_{18} & 0 & x_{19} & 0 & 0 & 0 & x_{20} & x_{21} & x_{22} \\ x_{23} & 0 & x_{24} & 0 & 0 & x_{25} & x_{26} & 0 & 0 & 0 \\ x_{27} & x_{28} & 0 & 0 & x_{29} & 0 & 0 & 0 & 0 & 0 \\ x_{30} & x_{31} & 0 & 0 & 0 & 0 & x_{32} & x_{33} & x_{34} & 0 \end{pmatrix}.$$

Proposition P is realizable $\iff \mathcal{V}_+(I_P) \neq \emptyset$.

Application 1: Realizability

Steinitz problem *Check whether an abstract polytopal complex is the boundary of an actual polytope.*

[Altshuler, Steinberg, 1985]: 4-polytopes and 3-spheres with 8 vertices.

The smallest non-polytopal 3-sphere has vertex-facet non-incidence matrix

$$S_P(\mathbf{x}) = \begin{pmatrix} 0 & 0 & 0 & 0 & 0 & x_1 & x_2 & x_3 & x_4 & x_5 \\ 0 & 0 & 0 & 0 & x_6 & x_7 & 0 & 0 & x_8 & x_9 \\ 0 & 0 & x_{10} & x_{11} & x_{12} & 0 & 0 & 0 & 0 & x_{13} \\ 0 & 0 & x_{14} & x_{15} & 0 & 0 & x_{16} & x_{17} & 0 & 0 \\ 0 & x_{18} & 0 & x_{19} & 0 & 0 & 0 & x_{20} & x_{21} & x_{22} \\ x_{23} & 0 & x_{24} & 0 & 0 & x_{25} & x_{26} & 0 & 0 & 0 \\ x_{27} & x_{28} & 0 & 0 & x_{29} & 0 & 0 & 0 & 0 & 0 \\ x_{30} & x_{31} & 0 & 0 & 0 & 0 & x_{32} & x_{33} & x_{34} & 0 \end{pmatrix}.$$

Proposition P is realizable $\iff \mathcal{V}_+(I_P) \neq \emptyset$.

In this case, $I_P = \langle 1 \rangle \Rightarrow$ no rank 5 matrix with this support \Rightarrow no polytope with the given facial structure.

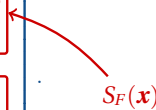
Application 2: Prescribability of faces

Lemma F face of $P \Rightarrow S_F$ submatrix of S_P and $I_F \subset I_P \cap \mathbb{C}[\mathbf{x}_F]$.
 F prescribable in $P \Leftrightarrow \mathcal{V}_+(I_F) = \mathcal{V}_+(I_P \cap \mathbb{C}[\mathbf{x}_F])$.

Application 2: Prescribability of faces

Lemma F face of $P \Rightarrow S_F$ submatrix of S_P and $I_F \subset I_P \cap \mathbb{C}[\mathbf{x}_F]$.
 F prescribable in $P \Leftrightarrow \mathcal{V}_+(I_F) = \mathcal{V}_+(I_P \cap \mathbb{C}[\mathbf{x}_F])$.

[Barnette, 1987]: 4-dimensional prism over a square pyramid with a non-prescribable cubical facet F

$$S_P(\mathbf{x}) = \begin{pmatrix} \boxed{\begin{matrix} x_1 & 0 \\ x_4 & 0 \\ x_7 & 0 \\ x_{10} & 0 \end{matrix}} & 0 & \boxed{\begin{matrix} 0 & x_2 & x_3 & 0 \\ 0 & 0 & x_5 & x_6 \\ x_8 & 0 & 0 & x_9 \\ x_{11} & x_{12} & 0 & 0 \end{matrix}} \\ x_{13} & 0 & x_{14} & 0 & 0 & 0 & 0 \\ \boxed{\begin{matrix} 0 & x_{15} \\ 0 & x_{18} \\ 0 & x_{21} \\ 0 & x_{24} \end{matrix}} & 0 & \boxed{\begin{matrix} 0 & x_{16} & x_{17} & 0 \\ 0 & 0 & x_{19} & x_{20} \\ x_{22} & 0 & 0 & x_{23} \\ x_{25} & x_{26} & 0 & 0 \end{matrix}} \\ 0 & x_{27} & x_{28} & 0 & 0 & 0 & 0 \end{pmatrix}.$$


$S_F(\mathbf{x})$

Application 2: Prescribability of faces

Lemma F face of $P \Rightarrow S_F$ submatrix of S_P and $I_F \subset I_P \cap \mathbb{C}[\mathbf{x}_F]$.
 F prescribable in $P \Leftrightarrow \mathcal{V}_+(I_F) = \mathcal{V}_+(I_P \cap \mathbb{C}[\mathbf{x}_F])$.

[Barnette, 1987]: 4-dimensional prism over a square pyramid with a non-prescribable cubical facet F

$$S_P(\mathbf{x}) = \begin{pmatrix} \boxed{\begin{matrix} x_1 & 0 \\ x_4 & 0 \\ x_7 & 0 \\ x_{10} & 0 \end{matrix}} & 0 & \boxed{\begin{matrix} 0 & x_2 & x_3 & 0 \\ 0 & 0 & x_5 & x_6 \\ x_8 & 0 & 0 & x_9 \\ x_{11} & x_{12} & 0 & 0 \end{matrix}} \\ x_{13} & 0 & x_{14} & 0 & 0 & 0 & 0 \\ \boxed{\begin{matrix} 0 & x_{15} \\ 0 & x_{18} \\ 0 & x_{21} \\ 0 & x_{24} \end{matrix}} & 0 & \boxed{\begin{matrix} 0 & x_{16} & x_{17} & 0 \\ 0 & 0 & x_{19} & x_{20} \\ x_{22} & 0 & 0 & x_{23} \\ x_{25} & x_{26} & 0 & 0 \end{matrix}} \\ 0 & x_{27} & x_{28} & 0 & 0 & 0 & 0 \end{pmatrix}. \quad \begin{matrix} \nearrow \\ S_F(\mathbf{x}) \end{matrix}$$

$$\dim(I_P \cap \mathbb{C}[\mathbf{x}_F]) = 15, \dim(I_F) = 16 \Rightarrow I_F \neq I_P \cap \mathbb{C}[\mathbf{x}_F]$$

Application 2: Prescribability of faces

Lemma F face of $P \Rightarrow S_F$ submatrix of S_P and $I_F \subset I_P \cap \mathbb{C}[\mathbf{x}_F]$.
 F prescribable in $P \Leftrightarrow \mathcal{V}_+(I_F) = \mathcal{V}_+(I_P \cap \mathbb{C}[\mathbf{x}_F])$.

[Barnette, 1987]: 4-dimensional prism over a square pyramid with a non-prescribable cubical facet F

$$S_P(\mathbf{x}) = \begin{pmatrix} \boxed{x_1} & \boxed{0} & 0 & \boxed{0} & \boxed{x_2} & \boxed{x_3} & \boxed{0} \\ \boxed{x_4} & \boxed{0} & 0 & \boxed{0} & \boxed{0} & \boxed{x_5} & \boxed{x_6} \\ \boxed{x_7} & \boxed{0} & 0 & \boxed{x_8} & \boxed{0} & \boxed{0} & \boxed{x_9} \\ \boxed{x_{10}} & \boxed{0} & 0 & \boxed{x_{11}} & \boxed{x_{12}} & \boxed{0} & \boxed{0} \\ x_{13} & 0 & x_{14} & 0 & 0 & 0 & 0 \\ \boxed{0} & \boxed{x_{15}} & 0 & \boxed{0} & \boxed{x_{16}} & \boxed{x_{17}} & \boxed{0} \\ \boxed{0} & \boxed{x_{18}} & 0 & \boxed{0} & \boxed{0} & \boxed{x_{19}} & \boxed{x_{20}} \\ \boxed{0} & \boxed{x_{21}} & 0 & \boxed{x_{22}} & \boxed{0} & \boxed{0} & \boxed{x_{23}} \\ \boxed{0} & \boxed{x_{24}} & 0 & \boxed{x_{25}} & \boxed{x_{26}} & \boxed{0} & \boxed{0} \\ 0 & x_{27} & x_{28} & 0 & 0 & 0 & 0 \end{pmatrix}. \quad S_F(\mathbf{x})$$

$$\dim(I_P \cap \mathbb{C}[\mathbf{x}_F]) = 15, \dim(I_F) = 16 \Rightarrow I_F \neq I_P \cap \mathbb{C}[\mathbf{x}_F]$$

With a bit more of work we conclude that F cannot be freely prescribed in P .

Application 3: Rationality

A combinatorial polytope is *rational* if it has a realization in which all vertices have rational coordinates.

Application 3: Rationality

A combinatorial polytope is *rational* if it has a realization in which all vertices have rational coordinates.

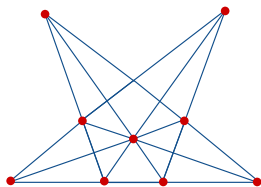
Lemma *A polytope P is rational $\Leftrightarrow \mathcal{V}_+(I_P)$ has a rational point.*

Application 3: Rationality

A combinatorial polytope is *rational* if it has a realization in which all vertices have rational coordinates.

Lemma *A polytope P is rational $\Leftrightarrow \mathcal{V}_+(I_P)$ has a rational point.*

We consider the following point-line arrangement in the plane [Grünbaum, 1967]:



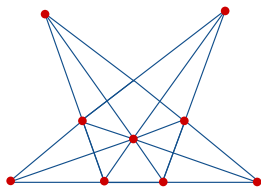
$$S_P(\mathbf{x}) = \begin{pmatrix} x_1 & 0 & x_2 & 0 & x_3 & x_4 & x_5 & x_6 & 0 \\ x_7 & x_8 & x_9 & 0 & x_{10} & 0 & 0 & x_{11} & x_{12} \\ x_{13} & x_{14} & 0 & x_{15} & x_{16} & x_{17} & x_{18} & 0 & 0 \\ x_{19} & x_{20} & 0 & x_{21} & 0 & 0 & x_{22} & x_{23} & x_{24} \\ x_{25} & 0 & x_{26} & x_{27} & 0 & x_{28} & 0 & 0 & x_{29} \\ 0 & 0 & x_{30} & x_{31} & x_{32} & 0 & x_{33} & x_{34} & x_{35} \\ 0 & x_{36} & 0 & x_{37} & x_{38} & x_{39} & 0 & x_{40} & x_{41} \\ 0 & x_{42} & x_{43} & 0 & x_{44} & x_{45} & x_{46} & 0 & x_{47} \\ 0 & x_{48} & x_{49} & x_{50} & 0 & x_{51} & x_{52} & x_{53} & 0 \end{pmatrix}$$

Application 3: Rationality

A combinatorial polytope is *rational* if it has a realization in which all vertices have rational coordinates.

Lemma A polytope P is rational $\Leftrightarrow \mathcal{V}_+(I_P)$ has a rational point.

We consider the following point-line arrangement in the plane [Grünbaum, 1967]:



$$S_P(\mathbf{x}) = \begin{pmatrix} x_1 & 0 & x_2 & 0 & x_3 & x_4 & x_5 & x_6 & 0 \\ x_7 & x_8 & x_9 & 0 & x_{10} & 0 & 0 & x_{11} & x_{12} \\ x_{13} & x_{14} & 0 & x_{15} & x_{16} & x_{17} & x_{18} & 0 & 0 \\ x_{19} & x_{20} & 0 & x_{21} & 0 & 0 & x_{22} & x_{23} & x_{24} \\ x_{25} & 0 & x_{26} & x_{27} & 0 & x_{28} & 0 & 0 & x_{29} \\ 0 & 0 & x_{30} & x_{31} & x_{32} & 0 & x_{33} & x_{34} & x_{35} \\ 0 & x_{36} & 0 & x_{37} & x_{38} & x_{39} & 0 & x_{40} & x_{41} \\ 0 & x_{42} & x_{43} & 0 & x_{44} & x_{45} & x_{46} & 0 & x_{47} \\ 0 & x_{48} & x_{49} & x_{50} & 0 & x_{51} & x_{52} & x_{53} & 0 \end{pmatrix}$$

Scaling rows and columns to set some variables to 1 (this does not affect rationality):

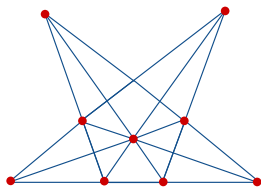
$$x_{46}^2 + x_{46} - 1 \in I_P$$

Application 3: Rationality

A combinatorial polytope is *rational* if it has a realization in which all vertices have rational coordinates.

Lemma A polytope P is rational $\Leftrightarrow \mathcal{V}_+(I_P)$ has a rational point.

We consider the following point-line arrangement in the plane [Grünbaum, 1967]:



$$S_P(\mathbf{x}) = \begin{pmatrix} x_1 & 0 & x_2 & 0 & x_3 & x_4 & x_5 & x_6 & 0 \\ x_7 & x_8 & x_9 & 0 & x_{10} & 0 & 0 & x_{11} & x_{12} \\ x_{13} & x_{14} & 0 & x_{15} & x_{16} & x_{17} & x_{18} & 0 & 0 \\ x_{19} & x_{20} & 0 & x_{21} & 0 & 0 & x_{22} & x_{23} & x_{24} \\ x_{25} & 0 & x_{26} & x_{27} & 0 & x_{28} & 0 & 0 & x_{29} \\ 0 & 0 & x_{30} & x_{31} & x_{32} & 0 & x_{33} & x_{34} & x_{35} \\ 0 & x_{36} & 0 & x_{37} & x_{38} & x_{39} & 0 & x_{40} & x_{41} \\ 0 & x_{42} & x_{43} & 0 & x_{44} & x_{45} & x_{46} & 0 & x_{47} \\ 0 & x_{48} & x_{49} & x_{50} & 0 & x_{51} & x_{52} & x_{53} & 0 \end{pmatrix}$$

Scaling rows and columns to set some variables to 1 (this does not affect rationality):

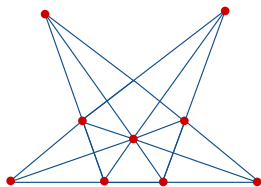
$$x_{46}^2 + x_{46} - 1 \in I_P \Rightarrow x_{46} = \frac{-1 \pm \sqrt{5}}{2} \Rightarrow \text{no rational realizations}$$

Application 3: Rationality

A combinatorial polytope is **rational** if it has a realization in which all vertices have rational coordinates.

Lemma A polytope P is rational $\Leftrightarrow \mathcal{V}_+(I_P)$ has a rational point.

We consider the following point-line arrangement in the plane [Grünbaum, 1967]:



$$S_P(\mathbf{x}) = \begin{pmatrix} x_1 & 0 & x_2 & 0 & x_3 & x_4 & x_5 & x_6 & 0 \\ x_7 & x_8 & x_9 & 0 & x_{10} & 0 & 0 & x_{11} & x_{12} \\ x_{13} & x_{14} & 0 & x_{15} & x_{16} & x_{17} & x_{18} & 0 & 0 \\ x_{19} & x_{20} & 0 & x_{21} & 0 & 0 & x_{22} & x_{23} & x_{24} \\ x_{25} & 0 & x_{26} & x_{27} & 0 & x_{28} & 0 & 0 & x_{29} \\ 0 & 0 & x_{30} & x_{31} & x_{32} & 0 & x_{33} & x_{34} & x_{35} \\ 0 & x_{36} & 0 & x_{37} & x_{38} & x_{39} & 0 & x_{40} & x_{41} \\ 0 & x_{42} & x_{43} & 0 & x_{44} & x_{45} & x_{46} & 0 & x_{47} \\ 0 & x_{48} & x_{49} & x_{50} & 0 & x_{51} & x_{52} & x_{53} & 0 \end{pmatrix}$$

Scaling rows and columns to set some variables to 1 (this does not affect rationality):

$$x_{46}^2 + x_{46} - 1 \in I_P \Rightarrow x_{46} = \frac{-1 \pm \sqrt{5}}{2} \Rightarrow \text{no rational realizations}$$

What else can we determine from I_P ?

A four-dimensional example

Let $P = \text{conv}\{0, 2e_1, 2e_2, 2e_3, e_{12} - e_3, e_4, e_{34}\}$, f -vector $(7, 17, 17, 7)$

$$S_P(\mathbf{x}) = \begin{pmatrix} 0 & x_1 & 0 & 0 & 0 & x_2 & 0 \\ x_3 & 0 & 0 & 0 & 0 & x_4 & 0 \\ x_5 & 0 & x_6 & 0 & 0 & 0 & x_7 \\ 0 & x_8 & x_9 & 0 & 0 & 0 & x_{10} \\ 0 & 0 & 0 & 0 & x_{11} & 0 & x_{12} \\ 0 & 0 & 0 & x_{13} & x_{14} & x_{15} & 0 \\ 0 & 0 & x_{16} & x_{17} & 0 & 0 & 0 \end{pmatrix}$$

There are 49 6-minors of $S_P(\mathbf{x})$, all binomials except 4 like:

$$x_4 x_5 x_{10} x_{11} x_{13} x_{16} - x_4 x_5 x_9 x_{12} x_{14} x_{17} + x_3 x_7 x_9 x_{11} x_{15} x_{17} - x_3 x_6 x_{10} x_{11} x_{15} x_{17}$$

After saturating the ideal of minors, the slack ideal is

$$I_P = \left\langle \begin{array}{ll} x_7 x_9 - x_6 x_{10}, & x_{10} x_{11} x_{13} x_{16} - x_9 x_{12} x_{14} x_{17}, \\ x_7 x_{11} x_{13} x_{16} - x_6 x_{12} x_{14} x_{17}, & x_2 x_8 x_{13} x_{16} - x_1 x_9 x_{15} x_{17}, \\ x_4 x_5 x_{13} x_{16} - x_3 x_6 x_{15} x_{17}, & x_2 x_8 x_{12} x_{14} - x_1 x_{10} x_{11} x_{15}, \\ x_4 x_5 x_{12} x_{14} - x_3 x_7 x_{11} x_{15}, & x_2 x_3 x_7 x_8 - x_1 x_4 x_5 x_{10}, \\ x_2 x_3 x_6 x_8 - x_1 x_4 x_5 x_9 & \end{array} \right\rangle$$

How nice can I_P be?

I_P cannot contain monomials

How nice can I_P be?

I_P cannot contain monomials

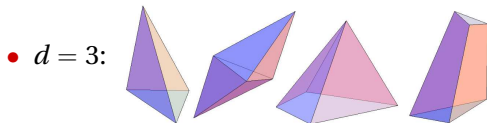
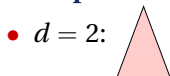
I_P can be generated by binomials, but usually not

How nice can I_P be?

I_P **cannot** contain monomials

I_P **can** be generated by binomials, but usually not

Examples:



• $d = 4$: 11 combinatorial classes of
PSD-minimal polytopes [GPRT17]

• Products of simplices

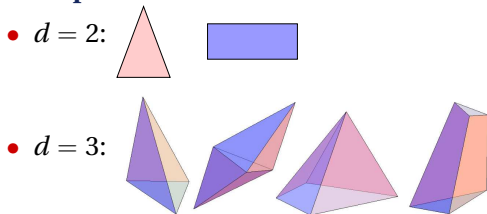
How nice can I_P be?

I_P **cannot** contain monomials

I_P **can** be generated by binomials, but usually not

I_P **can** be prime \rightarrow so can be toric!

Examples:



- $d = 4$: 11 combinatorial classes of PSD-minimal polytopes [GPRT17]

- Products of simplices

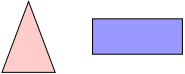
How nice can I_P be?

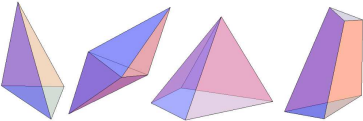
I_P **cannot** contain monomials

I_P **can** be generated by binomials, but usually not

I_P **can** be prime \rightarrow so can be toric!

Examples:

• $d = 2$: 

• $d = 3$: 

• $d = 4$: 11 combinatorial classes of
PSD-minimal polytopes [GPRT17]

• Products of simplices

All are toric and
projectively unique

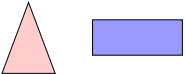
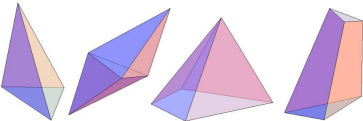
How nice can I_P be?

I_P **cannot** contain monomials

I_P **can** be generated by binomials, but usually not

I_P **can** be prime \rightarrow so can be toric!

Examples:

- $d = 2$: 
- $d = 3$: 
- $d = 4$: 11 combinatorial classes of PSD-minimal polytopes [GPRT17]
- Products of simplices

**Always
true?**

**All are toric and
projectively unique**

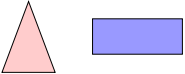
How nice can I_P be?

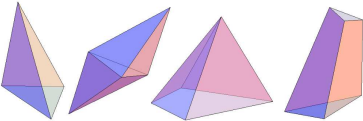
I_P **cannot** contain monomials

I_P **can** be generated by binomials, but usually not

I_P **can** be prime \rightarrow so can be toric!

Examples:

• $d = 2$: 

• $d = 3$: 

• $d = 4$: 11 combinatorial classes of
PSD-minimal polytopes [GPRT17]

• Products of simplices

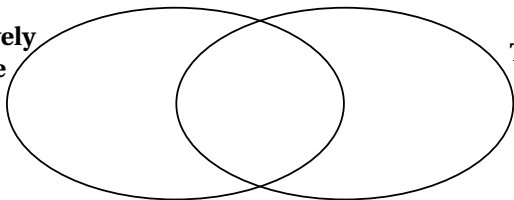
**Always
true?**

**All are toric and
projectively unique**

Is I_P the defining ideal of $\mathcal{V}_+(I_P)$? If I_P is toric, yes.

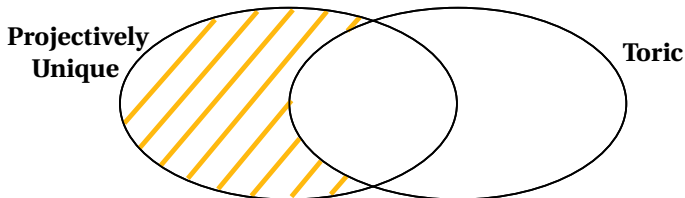
NO

**Projectively
Unique**



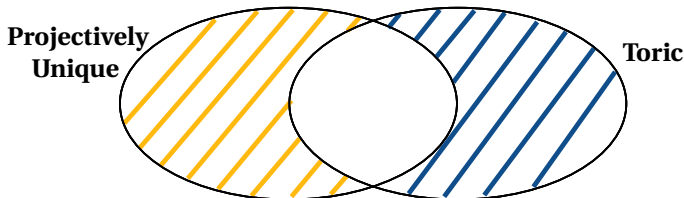
Toric

NO



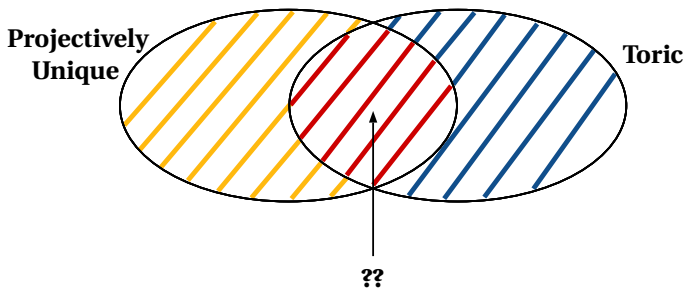
- In high enough dimension there are infinitely many projectively unique polytopes [Adiprasito, Ziegler, 2015], but only finitely many toric slack ideals in any dimension [Gouveia, Pashkovich, Robinson, Thomas, 2017].

NO



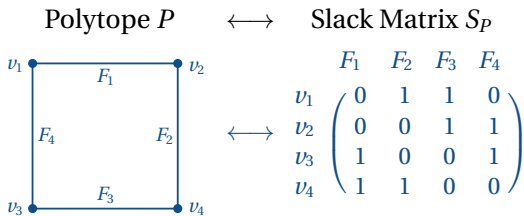
- In high enough dimension there are infinitely many projectively unique polytopes [Adiprasito, Ziegler, 2015], but only finitely many toric slack ideals in any dimension [Gouveia, Pashkovich, Robinson, Thomas, 2017].
- In dimension 5 there are non-projectively unique polytopes that have toric slack ideals [Gouveia, M, Thomas, Wiebe, 2017].

NO

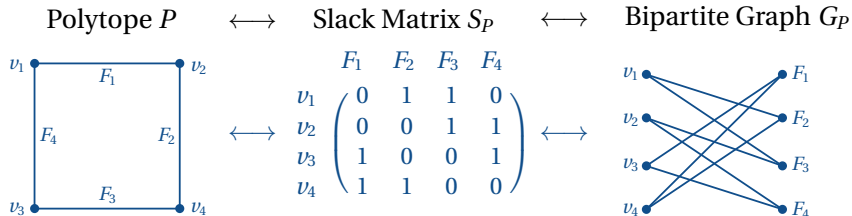


- In high enough dimension there are infinitely many projectively unique polytopes [Adiprasito, Ziegler, 2015], but only finitely many toric slack ideals in any dimension [Gouveia, Pashkovich, Robinson, Thomas, 2017].
- In dimension 5 there are non-projectively unique polytopes that have toric slack ideals [Gouveia, M, Thomas, Wiebe, 2017].
- *Which polytopes are toric and PU?*

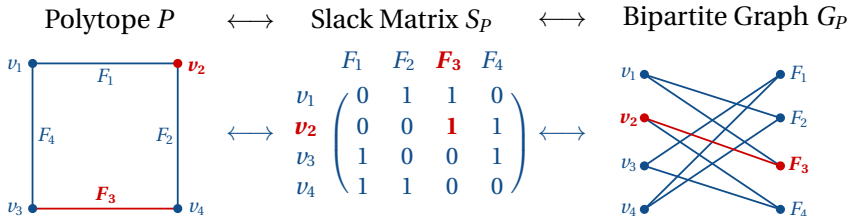
Non-incidence graph of a polytope



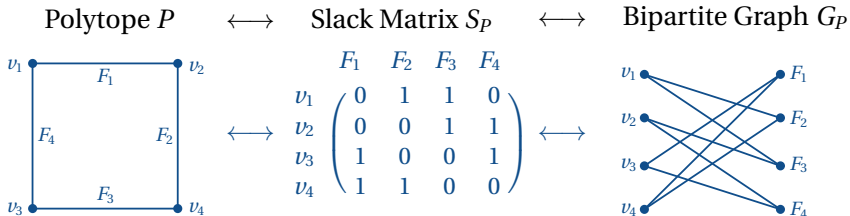
Non-incidence graph of a polytope



Non-incidence graph of a polytope



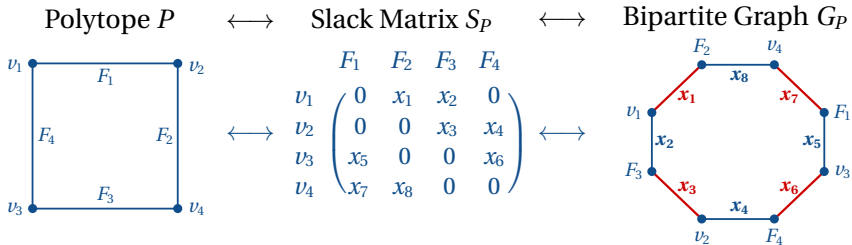
Non-incidence graph of a polytope



Toric ideal of a graph is a well-studied object:

$$T_P = \langle \mathbf{x}^{C^+} - \mathbf{x}^{C^-} : C \text{ cycles in } G_P \rangle$$

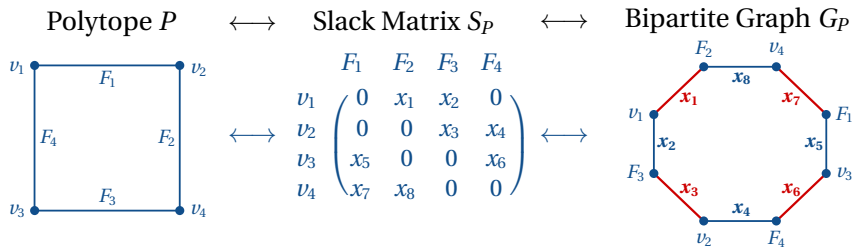
Non-incidence graph of a polytope



Toric ideal of a graph is a well-studied object:

$$T_P = \langle \mathbf{x}^{C^+} - \mathbf{x}^{C^-} : C \text{ cycles in } G_P \rangle \quad T_P = \langle x_2 x_4 x_5 x_8 - x_1 x_3 x_6 x_7 \rangle$$

Non-incidence graph of a polytope



Toric ideal of a graph is a well-studied object:

$$T_P = \langle \mathbf{x}^{C^+} - \mathbf{x}^{C^-} : C \text{ cycles in } G_P \rangle \quad T_P = \langle x_2 x_4 x_5 x_8 - x_1 x_3 x_6 x_7 \rangle$$

What is the relation between I_P and T_P ?

Graphic Polytopes

Geometric meaning of polytopes for which $I_P \subseteq T_P$.

Graphic Polytopes

Geometric meaning of polytopes for which $I_P \subseteq T_P$.

A polytope P is **2-level** if it has a slack matrix in which every positive entry is one, i.e., $S_P(\mathbb{1})$ is a slack matrix of P .

Graphic Polytopes

Geometric meaning of polytopes for which $I_P \subseteq T_P$.

A polytope P is **2-level** if it has a slack matrix in which every positive entry is one, i.e., $S_P(\mathbb{1})$ is a slack matrix of P .

Definition

- A polytope P is **morally 2-level** if $S_P(\mathbb{1})$ lies in the slack variety of P .
- I_P **graphic** if $I_P = T_P$.

Graphic Polytopes

Geometric meaning of polytopes for which $I_P \subseteq T_P$.

A polytope P is **2-level** if it has a slack matrix in which every positive entry is one, i.e., $S_P(\mathbb{1})$ is a slack matrix of P .

Definition

- A polytope P is **morally 2-level** if $S_P(\mathbb{1})$ lies in the slack variety of P .
- I_P **graphic** if $I_P = T_P$.

Theorem (Gouveia, M, Thomas, Wiebe, 2017)

- 1 A polytope P is *morally 2-level* $\Leftrightarrow I_P \subseteq T_P$.
- 2 I_P is *graphic* $\Leftrightarrow I_P$ toric and P projectively unique.

Conclusions

- Slack matrix encodes the combinatorics of polytopes
- Positive part of the slack variety as a model of the realization space for modding out projective equivalence
- New characterization of class of projectively unique polytopes via slack ideal: graphic polytopes are PU

Conclusions

- Slack matrix encodes the combinatorics of polytopes
- Positive part of the slack variety as a model of the realization space for modding out projective equivalence
- New characterization of class of projectively unique polytopes via slack ideal: graphic polytopes are PU

What's next?

- Continue to improve this new dictionary between algebra and combinatorics of polytopes
- Is the slack ideal prime or radical?
- Classify polytopes with toric slack ideal
- Polytopes with binomial non-toric slack ideal?
- New classes of projective unique polytopes

Thank you!